



09-21-00

A

LAW OFFICES  
BARRY R. LIPSITZ

BRADFORD GREEN, BUILDING 8  
755 MAIN STREET  
MONROE, CONNECTICUT 06468

PATENTS, TRADEMARKS, COPYRIGHTS

BARRY R. LIPSITZ  
RALPH F. HOPPIN  
DOUGLAS M. McALLISTER

TELEPHONE: (203)459-0200  
FACSIMILE: (203)459-0201

**BOX PATENT APPLICATION**

Commissioner for Patents  
Washington, D.C. 20231

**UTILITY PATENT APPLICATION TRANSMITTAL  
FILED UNDER 37 C.F.R. §1.53(b)**

jc564 U.S. PTO  
09/666901  
09/20/00

Sir:

Transmitted herewith for filing is the patent application of:

Inventor(s): **Vincent LIU, Siu-Wai WU, Michael CASTELOES, Robert J. STONE and Rebecca LAM**

Title: **METHOD AND APPARATUS FOR DETERMINING A TRANSMISSION BIT RATE IN A STATISTICAL MULTIPLEXER**

**APPLICATION ELEMENTS:**

(1)  Patent Application Specification, including Abstract and claims - 37 pages  
(2)  Seven (7) sheets of formal drawings, together with transmittal letter  
(3)  A check in the amount of \$730.00 to cover the [ X ] filing fee (\$690.00 ) and/or [ X ] Assignment Recordal Fee (\$40) is enclosed.

[ ] Before calculating the fee, cancel claim(s)  
[ ] Before calculating the fee, see copy of Preliminary Amendment filed in parent application \_\_\_\_\_ (attached hereto.)

Basic Fee							\$ 690.00
Multiple Dependent Claims (\$260.)							
Foreign Language Surcharge (\$130.)							
	For	No. Filed	-	No. Extra		Rate	
EXTRA	Total Claims	13	20	0	x	\$18.	= \$ 0.00
CLAIMS	Independent Claims	2	3	0	x	\$78.	= \$ 0.00
<b>TOTAL FILING FEE</b>							= \$ 690.00

The Commissioner is hereby authorized to charge any deficiency in the payment of the required fee(s) or credit any overpayment to Deposit Account No. 50-0625.

**Box Patent Application  
Commissioner for Patents  
Page 2 of 3**

(4) [X] Declaration and Power of Attorney form -- 3 pages

- [X] Newly executed
- [ ] Copy from a prior application
- [ ] Deletion of inventor(s) – signed statement attached deleting inventor(s) named in the prior application.

(5) [ ] Small Entity Declaration

- [ ] Newly executed
- [ ] Copy from a prior application. Status still proper and desired.

**ACCOMPANYING APPLICATION PARTS:**

(6) [X] Assignment document

- [X] Newly executed (with \$40.00 recordal fee) and separate transmittal Form PTO-1595
- [ ] Copy from a prior application

(7) [ ] Preliminary Amendment

(8) [ ] Certified Copy of Priority Document, together with separate transmittal letter

(9) [ ] Information Disclosure Statement, together with PTO Form 1449 and copies of cited references

(10) [X] Return receipt postage prepaid postcard

(11) [X] Express Mail Certificate (Mailing Label No. EL 632260048 US)

(12) [ ] Other:

- [ ]

(13) [ ] If a **CONTINUING APPLICATION**, check appropriate box, and supply the requisite information below:

Continuation     Divisional     Continuation-in-part (CIP)

of prior application no.:

Prior application information:

Examiner: \_\_\_\_\_ Group/Art Unit: \_\_\_\_\_

Status: \_\_\_\_\_

**Box Patent Application**  
**Commissioner for Patents**  
**Page 3 of 3**

FOR CONTINUATION or DIVISIONAL APPLICATIONS only: The entire disclosure of the prior application, from which an oath or declaration is supplied under paragraph 4(b) above, is considered a part of the disclosure of the accompanying continuation or divisional application and is hereby incorporated by reference. The incorporation can only be relied upon when a portion has been inadvertently omitted from the submitted application parts.

(14) [ ] Please amend the specification by inserting before the first line the sentence:

(15) [ ] This application claims the benefit of \_\_\_\_\_ Patent Application  
No. \_\_\_\_\_ filed on \_\_\_\_\_ .

(16) [ X ] Correspondence address:  
**Barry R. Lipsitz**  
**Law Offices of Barry R. Lipsitz**  
**755 Main Street, Building 8**  
**Monroe, CT 06468**  
**Telephone: (203) 459-0200**

Respectfully submitted,



---

Barry R. Lipsitz  
Attorney for Applicant(s)  
Registration No. 28,637  
755 Main Street  
Monroe, CT 06468  
(203) 459-0200

**ATTORNEY DOCKET NO.: GIC-620**  
**Date: September 20, 2000**

PATENT

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Application of: )  
Liu et al. )  
Application No.: )  
Filed: Herewith )  
For: **METHOD AND APPARATUS FOR DETERMINING A TRANSMISSION BIT RATE IN A  
STATISTICAL MULTIPLEXER**

**BOX PATENT APPLICATION**

Commissioner for Patents  
Washington, D.C. 20231

**EXPRESS MAIL CERTIFICATE**

"Express Mail" mailing label number: EL 632260048 US  
Date of Deposit: September 20, 2000

I hereby certify that the attached:

- Check in the amount of \$730.00 (Filing fee (\$690) and recordal fee (\$40));
- Return receipt postage prepaid postcard;
- Transmittal letter for new patent application;
- Patent Application Specification, including Abstract and Claims (37 pages);
- Seven (7) sheets of formal drawings, together with transmittal letter;
- Declaration and Power of Attorney form;
- Assignment, together with Form PTO-1595.

are being deposited with the United States Postal Service "Express Mail Post Office to Addressee" service under 37 C.F.R. 1.10 on the date indicated above and is addressed to: **BOX PATENT APPLICATION, Commissioner for Patents, Washington, D.C. 20231**

Michele Hollis  
(Typed or printed name of person mailing paper or fee)

  
(Signature of person mailing paper or fee)

Respectfully submitted,

  
Barry R. Lipsitz  
Attorney for Applicant(s)  
Registration No. 28,637  
755 Main Street  
Monroe, CT 06468  
(203) 459-0200

Date: **September 20, 2000**  
ATTORNEY DOCKET NO.: **GIC-620**

METHOD AND APPARATUS FOR DETERMINING A TRANSMISSION BIT  
RATE IN A STATISTICAL MULTIPLEXER

BACKGROUND OF THE INVENTION

The present invention relates to a statistical  
5 multiplexer for coding and multiplexing multiple  
channels of digital television data.

Digital television has become increasingly popular  
due to the high quality video image it provides, along  
with informational and entertainment features, such as  
10 pay-per-view, electronic program guides, video-on-  
demand, stock, weather and stock information, Internet  
hyperlinks, and so forth. Such television data can be  
communicated to a user, for example, via a broadband  
communication network, such as a satellite or cable  
15 television network, or via a computer network.

However, due to the bandwidth limitations of the  
communication channel, it is necessary to adjust a bit  
rate of the digital video programs that are encoded and  
multiplexed for transmission in a compressed bit  
20 stream. A goal of such bit rate adjustment is to meet  
the constraint on the total bit rate of the multiplexed  
stream, while also maintaining a satisfactory video  
quality for each program.

Accordingly, various types of statistical  
25 multiplexers have been developed that evaluate

statistical information of the source video that is being encoded, and allocate bits for coding the different video channels accordingly. For example, video channels that have hard-to-compress video, such as a fast motion scene, can be allocated more bits, while channels with relatively easy to compress scenes, such as scenes with little motion, can be allocated fewer bits.

However, there is a need for an improved statistical multiplexing system. Such a system should employ a number of individual encoders that encode data from a number of incoming channels of source video data. This data may be obtained from a storage media, live feed, or the like.

The system should dynamically allocate bits to the individual encoders to encode frames of video data from the channels.

The system should distinguish between an encoding bit rate allocation and a transmission bit rate allocation.

The system should assign a transmission bit rate to each video channel to prevent a mismatch between the input and output bits of a modeled decoder buffer.

Moreover, the system should check for impending decoder buffer overflow or underflow events to set minimum and maximum limits on the transmission bit rate. These limits should be set based on whether a new transmission bit rate can be implemented before the

decoding time stamp (DTS) of the current or next frame.

The system should be usable with essentially any type of video data, including high-definition (HD) and standard-definition (SD) television (TV).

5 The present invention provides a system having the above and other advantages.

## SUMMARY OF THE INVENTION

The present invention relates to a statistical multiplexer for coding and multiplexing multiple channels of digital television data.

5 Bandwidth is dynamically allocated among a number of variable bit rate (VBR) video services that are multiplexed to form a fixed bit rate transport bit stream.

10 In particular, a transmission bit rate is calculated by delaying an encoding bit rate by an amount equal to a system delay. An upper bound (max) and lower bound (min) are placed on the transmission bit rate to protect the decoder buffer.

15 The stat mux system includes three distinct parts:  
1) The collection of visual characteristics and complexity information for individual video channels and a need parameter is generated for each video channel to indicate how difficult it is to compress that channel. This process is repeated once per frame  
20 and it is done by the individual single-channel encoders (which could be SD and/or HD).

25 2) The most up-to-date need parameters from all the video channels are collected by a quantization level processor (QLP), or rate control processor. The rate control processor assigns an encoding bandwidth to each video channel in the form of an encoding bit rate. Each channel receives a different encoding bit rate.

based on its need parameter in relation to the needs of all the other channels. The encoding bit rate is used to control the video encoding of individual channels. The rate control processor also assigns transmission 5 bit rates to the channels, which determine how many bits are sent by each video channel to a decoder.

3) The single-channel encoder uses the encoding bit rate it is given to perform video compression. The primary task here is a rate control function, which involves using the encoding bit rate and the relative complexities of different frame-types (i.e., I, B and P types) to assign a target bit budget for each frame it is about to encode.

A particular method for processing a plurality of channels in a statistical multiplexer includes the steps of: allocating an encoding bit rate for coding a current picture of each channel according to a bit rate need parameter thereof, and allocating a transmission bit rate for transmitting the current picture of each channel after encoding thereof. For each channel, the transmission bit rate is allocated in accordance with the channel's encoding bit rate after the encoding bit rate has been allocated to the channel, and following a system delay of a modeled buffer of a decoder, to 20 minimize a rate mismatch between an input and an output of the modeled decoder buffer.

Moreover, min. and max. limits are provided for the transmission bit rate to avoid a potential overflow

or underflow of a modeled decoder buffer. These limits take into account a delay in implementing a new bit rate, and also account for whether the new bit rate can be implemented before the decoding time of the current picture, the next picture, and/or the next, next picture.

Note that the pictures can be, e.g., frames or fields.

A corresponding apparatus is also presented.

#### BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 illustrates a statistically multiplexed multi-channel encoding system for use in accordance with the present invention.

5 FIG. 2 illustrates an encoder for standard definition television data for use in accordance with the present invention.

10 FIG. 3 illustrates an encoder for high-definition television data for use in accordance with the present invention.

15 FIG. 4(a) illustrates a time line where a new maximum transmission bit rate can be implemented before the decoding time stamp (DTS) of the next frame, in accordance with the present invention.

FIG. 4(b) illustrates a time line where a new maximum transmission bit rate can not be implemented before the decoding time stamp (DTS) of the next frame, in accordance with the present invention.

20 FIG. 4(c) illustrates a time line where a new minimum transmission bit rate can be implemented before the decoding time stamp (DTS) of the current frame, in accordance with the present invention.

25 FIG. 4(d) illustrates a time line where a new minimum transmission bit rate can not be implemented before the decoding time stamp (DTS) of the current frame, in accordance with the present invention.

FIG. 5 illustrates a method for setting minimum

and maximum limits on a transmission bit rate in accordance with the present invention.

## DETAILED DESCRIPTION OF THE INVENTION

The present invention relates to a statistical multiplexer for coding and multiplexing multiple channels of digital television data.

FIG. 1 illustrates a statistically multiplexed multi-channel encoding system for use in accordance with the present invention.

The encoding system 100 includes L buffer/need parameter calculation functions 102, 104, ..., 106 that receive corresponding uncompressed source video inputs. The functions 102, 104, ..., 106 provide the need parameter data to a rate control processor 125, which in turn provides a corresponding encoding bit rate allocation to each of the encoders 112, 114, ..., 116. The encoders may provide feedback information to the rate control processor regarding the actual encoding bit rate. The encoded data is provided to a mux 120 to provide a multiplexed bitstream, then to a transport packet buffer 130, and to a transmitter 135 for transmission across a channel.

The rate control processor 125 may receive a fullness signal from the transport packet buffer 130.

At a decoding side 180, a receiver 182, decoder buffer 184, demux 186, and decoder 188 are provided to output a decoded video signal, e.g., for display on a television.

FIG. 2 illustrates an encoder for standard

definition television for use in accordance with the present invention.

The encoder 112, which is an example one of the encoders 112, 114, ..., 116 of FIG. 1, encodes a single channel of input data, and includes a compressor 210 that performs conventional data compression, including motion compensation (for P- and B-frames), discrete cosine transform (DCT) and quantization. A video first-in, first-out (FIFO) buffer 230 temporarily stores the compressed data, and a packet processor 250 forms packets of the compressed data with appropriate header information, e.g., according to the MPEG-2 or other video standard.

FIG. 3 illustrates an encoder for high-definition television for use in accordance with the present invention.

The encoder 300 encodes a single channel of input data. However, a panel splitter 305 divides up a video frame such that different sub-regions, or panels, of the frame are routed to respective different compressors 310-324. Eight compressors are shown as an example only. Typically, the same sub-region of successive frames is assigned to the same compressor.

A master compression controller (MCC) 370 controls the compression of the data at each compressor via a peripheral component interconnect (PCI) bus 325, and the compressed data is output to a video FIFO 330 for temporary storage. The compressed data is formed into

packets for transport at a packet processor 350.

A encoding bit rate need parameter for the HDTV channel is determined by the MCC 370 by summing a need parameter for each of the panel compressors. Other 5 statistical information, such as motion estimation scores and the like, are also summed from each compressor

Note that it is possible to combine both SDTV and HDTV encoders in a single stat mux group. In this case, the encoder 300 is an example one of the encoders 10 112, 114,..., 116 of FIG. 1. For example, one HDTV encoder may be combined with two or three SDTV encoders in a stat mux group.

#### Overview

15 A key part of a statistically multiplexed multi-channel encoding system of the invention is the calculation of the need parameter.

The visual characteristics and complexity information regarding the source video are collected 20 and condensed into a single parameter, which is referred to as the "need parameter". A need parameter is calculated for each video channel, and is updated once per frame whenever a new video frame is processed by the corresponding single-channel encoder 112, 114, 25 ..., 116. Optionally, the need parameter can be updated more often, such as multiple times per frame. Moreover, for field-picture mode, the need parameter can be updated once per field.

Discussion

In the following description of a stat mux, each video service is assumed to provide a picture complexity measure, such as an ME score or activity level, to the rate control processor 125, which handles the tasks of allocating bandwidth for each television service provider (TSP), e.g., channel, and modulating the transmission rates for each channel. In an encoder with look ahead capability, the ME score can be replaced by other measurements such as the actual number of bits coded under a constant quantization level (QL).

For the high-definition encoder that processes multiple panels of a frame in parallel, the encoders 112, 114, ..., 116 collect the ME scores from all the panels and compute the sum along with other parameters such as average pixel level (APL), picture resolution, frame rate, frame type (I, B or P) and total intra-frame activity. It also keeps a record of the sizes and average QL for past frames. Based on the information available, plus the look ahead parameters from scene change, fade and film detection, the MCC 370 can derive a need parameter for that video channel.

As the rate control processor 125 receives an updated need parameter from a buffer/need parameter calculation function 102, 104, ..., 106, it reallocates the bandwidths for all the video services based on the latest information. The bandwidth allocation is sent

back to each encoder 112, 114, ..., 116 in the form of an encoding bit rate. Moreover, the rate control processor 125 uses the bandwidth allocation to compute bit budgets for encoding. It keeps an approximate 5 video buffering verifier (VBV) model, such as is known from the MPEG standard, to ensure that each frame is encoded within acceptable size limits.

Note that the VBV model is only approximate because the actual transmission rate changes that occur at the decode time of a frame cannot be precisely modeled in advance, at the time of encoding. The rate control processor 125 keeps a bit accurate model of the decoder buffer 184, and if it is given the sizes of each encoded frame along with the decoding time stamp (DTS), the min. and max. limits on the transmission rate can be calculated and used before a transmission rate change is issued. As known from the MPEG standard, a DTS is a field that is present in a PES packet header that indicates the time that an access 15 unit (e.g., picture) is decoded in a system target decoder.

Here, the min. and max. rates are used for decoder buffer protection. Under normal conditions, one expects transmission rate changes to follow (with a 25 time lag) the changes in the encoding rate as communicated by the encoders to the rate control processor. The time lag, or system delay, is the total buffer system delay through both encode and decode

buffers, and is assumed to be the same for all video services in the stat mux group. The system delay includes the time between encoding and decoding of a picture. In the case of multiplexing HD and SD channels, their system delays can be approximately the same if a relative weighting of, e.g., five-to-one between HD and SD channels. is assumed.

Since all the video services need not be frame-synchronized, the encoding bit rates and transmission rates are updated as frequently as the rate control processor can handle. Encoding bit rates are passed on (with the buffer system delay) to become the transmission rate for that channel.

1. Assignment Of Transmission Bit rate To Each Video Channel

At the heart of the statistical multiplexing system is the bandwidth allocation process, which is performed by the rate control processor 125. In our stat mux system, a distinction can be made between two different concepts of bandwidth: one is the allocation of bandwidth to each video channel for the purpose of encoding. From the perspective of the rate control processor, each single-channel video encoder is allocated an encoding bit rate (BR) which is used to determine the targeted size to encode a video frame. The single-channel encoder receives updates on its encoding bit rate every QL/BR period, which may occur, e.g., more than ninety times per second. However

before encoding a particular video frame, only the latest encoding bit rate is used to calculate the targeted budget size for encoding that frame.

The other bandwidth allocation process in accordance with the invention involves the assignment of transmission bit rate to each video channel. This can also be calculated every QL/BR period, or at a faster or slower rate. The transmission rate for a video channel determines how many bits of compressed video should be sent by that channel during a QL/BR period. The sum of transmission bit rates over all the video channels for a particular QL/BR period should be equal to, or less than, the total video bandwidth given to the stat mux system. Similarly, the sum of encoding bit rates over all the video channels for a particular QL/BR period should be equal to, or less than, the total video bandwidth given to the system.

### 1.1 Relationship between Encoding and Transmission Bit Rates

For an MPEG-2 video encoding/decoding system, the concept of buffer system delay refers to the time lag from the time a video frame is first encoded at the encoder to the time that video frame is decoded by the decoder, e.g., as designated by the frame's DTS. For most MPEG-2 systems, the buffer system delay is constant once the encoding/decoding system is configured to a certain bit rate. A constant buffer system delay may also be assumed for the present stat

mux system.

Furthermore, the transmission bit rate is chosen to approximately be a time-lagged version of the encoding bit rate. The time lag between the two rates is equal to the buffer system delay.

The reason for this particular choice is as follows. If one ignores any latency in the transmission path between transport packet buffer 130 at the encoding side and the decoder buffer 184, the transmission bit rate is the same as the rate at which video bits arrive at the decoder buffer. Since the encoding bit rate is used in determining the budget for encoding a video frame, one can also think of it as the decoding rate when that particular frame is decoded.

The bit arrival rate at the decoder buffer (i.e., the transmission bit rate) equals the encoding bit rate delayed by the buffer system delay time, and the decoding bit rate also equals the encoding bit rate delayed by the buffer system delay time. Therefore, when a video frame is decoded, the decoding rate (which controls the output rate of the decode buffer) roughly equals the bit arrival rate at the input of the decoder buffer. Therefore, the present invention minimizes any possible rate mismatch between the output and the input to the decoder buffer. In a fixed rate system, this mismatch never occurs since the encoding bit rate and the transmission bit rate are the same fixed rate.

In the above discussion, the encoding/decoding bit

rate refers to the rate at which a video frame is encoded/decoded without considering the effects of different frame types. The assignment of bit budgets according to frame types is done by the rate control function at the individual single-channel encoder. For example, for a given encoding bit rate, the rate controller assigns a large bit budget to an I-frame if that is the type for the frame it is about to encode. Conversely, a B-frame is assigned a much smaller bit budget.

1.2 Min And Max Bit Rate Requests From The Packet Processor

Note that the transmission bit rate is only approximately a time-lagged version of the encoding bit rate. The rate control processor 125 has the freedom to set the transmission bit rate differently if necessary. When the transmission bit rate deviates from the encoding bit rate, the frames that are already encoded and stored in the transport packet buffer 130 may cause the decoder buffer 184 to overflow or underflow. This is because those frame were originally encoded under the assumption that the transmission rate follows the encoding bit rate exactly. The transmission rate may deviate from the encoding bit rate for several reasons, e.g., temporary congestion in the transmission network, or imprecision in the actual number of bits sent over a QL/BR period of time.

In the present stat mux system, such deviations

are allowed for while still maintaining overflow/underflow protections for the decode buffer. This is achieved by having the packet processor 250 or 350 check for impending overflow/underflow events, and  
5 communicates back to the rate control processor 125 through a set of maximum and minimum transmission bit rate requests. The maximum bit rate defines the highest rate at which the packet processor may send video packets for that channel without overflowing that channel's decoder buffer. The minimum bit rate request informs the rate control processor 125 of the lowest rate at which the individual video channel must send its video bits to get the complete video frame over to the decoder before the frame's decoding time starts.

10  
15 For the packet processor (250 or 350) to compute the min. and max. limits on the transmission bit rate, it needs to keep track of the decoder buffer through a decoder buffer model, including the current status of the buffer, the number of frames in it, the DTS and sizes of each frame, and the DTS and size of the frame  
20 that is currently being transmitted.

25 A picture is transmitted in advance of its DTS time, of course, e.g., up to 0.5 sec before. The picture data is stored in the decoder's buffer prior to being decoded at the DTS time. The DTS is based on a program clock reference (PCR), which is a clock that runs on an encoder board. This time is periodically sent from the encoder to the decoder to synchronize the

devices. The PCR clock runs at 27 MHz in a common implementation.

Before each new transmission rate is issued, the rate control (QL) processor 125 polls all the channels for min. and max. allowable bit rates. The packet processor for each channel calculates these limits as follows.

*CT* is the current time, *Next\_DTS* is the time for the decoder to decode its next frame, and the DTS after that is *Next\_next\_DTS*. "Latency", "delay" or "implementation delay" refer to the interval from the current time until the time when the bit rate can be implemented. This delay accounts for factors such as a signal travel time between the packet processor and the rate control processor, and a time for the rate control processor to calculate a new transmission bit rate.

The delay depends on the hardware configuration.

The transmission bit rate is updated in successive cycles, which can vary from the cycles for updating the encoding bit rate. Thirty-two transmission bit rate update cycles between DTSs can be used. The interval between DTSs is the inverse of the frame rate, e.g., 1/24 or 1/30 sec. The latency is typically a fraction (e.g., 1/3) of a transmission bit rate update cycle.

The *current\_br* is the transmission bit rate for the current cycle received from the rate control processor 125, and scaled by 184/188 to remove the transport packet header bits. The packet processor of

each channel outputs packets according to the transmission bit rate assigned to it. Assume an MPEG transport packet having 184 payload bytes and 4 header bytes is used. However, other transport schemes may be  
5 used, in which case this adjustment is modified as required.

In the following equations, *current\_br* is further adjusted by a "margin" of 61 kbytes/sec to account for PCR, PES headers and time stamps. For an HDTV encoder,  
10 the margin is doubled to 122 kbytes/sec. These values, again, are specific to the MPEG coding standard, but the invention is suitable for use with other coding and transmission schemes as well.

The maximum bit rate limit is used to prevent the modeled decoder buffer from overflowing for each  
15 channel. The requirement is that the decoder buffer should not overflow at the next DTS time. By monitoring the amount of space in the decoder buffer at the *Next\_DTS* time (*decode\_buffer\_size - decode\_buffer*),  
20 assuming no new bits are transmitted, and adjusting for the latency in the rate control processor's response, the maximum bit rate is given by the amount of empty buffer space divided by the time until the *Next\_DTS*.

The timeline 400 of FIG. 4(a) shows the case where  
25 the time for implementing the next transmission bit rate update is prior to the *Next\_DTS*. As shown in the timeline 420 of FIG. 4(b), if the time for implementing the next bit rate update is after the *Next\_DTS*, the

buffer overflow condition should be checked for the *Next\_next\_DTS*. The maximum bit rate calculation is expressed in the following C-language syntax.

If (*CT* + *latency*) < *Next\_DTS* then

5                $\max\_br = (\text{decode\_buffer\_size} - \text{decode\_buffer} - (\text{current\_br} + \text{margin}) * \text{latency}) / (\text{Next\_DTS} - \text{CT} - \text{latency});$

else

10               $\max\_br = (\text{decode\_buffer\_size} - \text{decode\_buffer} + \text{picture\_size\_of}(\text{picture\_at\_Next\_DTS}) - (\text{current\_br} + \text{margin}) * \text{latency}) / (\text{Next\_next\_DTS} - \text{CT} - \text{latency})$

When (*CT*+*latency*)<*Next\_DTS*, the numerator represents the modeled decoder buffer fullness at time (*CT*+*latency*), and the denominator represents the time period between (*CT*+*latency*) and (*Next\_DTS*). *Max\_br* is then the bit rate that will consume the remaining buffer space when implemented in the time period between (*CT*+*latency*) and (*Next\_DTS*). See FIG. 4(a).

15              When (*CT*+*latency*) $\geq$ *Next\_DTS*, the numerator represents the modeled decoder buffer fullness at time (*CT*+*latency*), and the denominator represents the time period between (*CT*+*latency*) and (*Next\_next\_DTS*). *Max\_br* is then the bit rate that will consume the remaining buffer space when implemented in the time period between (*CT*+*latency*) and (*Next\_next\_DTS*).

20              *picture\_size\_of(picture\_at\_Next\_DTS)* is the size (e.g., in bits) of the picture whose DTS is *Next\_DTS*. See FIG. 4(b).

In the code for min. bit rate calculations, below,

the term "current frame" refers to the coded video  
 frame that is currently being transmitted. The minimum  
 bit rate is derived from the requirement that the  
 current frame of compressed data must be completely  
 5 transmitted before the DTS for that frame, otherwise  
 the decoder buffer may underflow during the decoding of  
 that frame. First, check if the next bit rate change  
 occurs before the DTS time of the current frame ( i.e.  
 $CT + latency < DTS\_of\_current\_frame$ ). If so, as shown  
 10 in the timeline 440 of FIG. 4(c), it is meaningful to  
 calculate a  $min\_br$  defined as the minimum constant bit  
 rate required to send the remainder of the frame before  
 the DTS time. If  $min\_br$  turns out to be zero or less  
 (this could happen if, for a given latency, the  
 15  $no\_of\_bits\_left\_current\_frame \leq current\_br * latency$ )  
 then  $min\_br$  is set to zero. This means there is no  
 minimum requirement on the new bit rate.

For the special case when the DTS is so close to  
 the current time that  $CT + latency >$   
 20  $DTS\_of\_current\_frame$ , as shown in the timeline 460 of  
 FIG. 4(d),  $min\_br$  is not meaningful. In this case,  
 $min\_br$  is set to zero if the old bit rate can finish  
 delivering all the remaining bits before DTS occurs;  
 otherwise  $min\_br$  is set to a large enough value to  
 25 mitigate any potential decoder buffer underflow  
 problem. The minimum bit rate calculation is expressed  
 in the following C-language syntax.

If ( $latency < DTS\_of\_current\_frame - CT$ ) then

```

min_br = ( no_of_bits_left_current_frame - (current_br -
margin) * latency ) /
(DTS_of_current_frame - CT - latency );
If (min_br < 0) then min_br = 0;
5      else
        If ( No_of_bits_left_current_frame < (current_br - margin)
*(DTS_of_current_frame - CT) ) then
            min_br = 0;
        else
10          min_br = HARDWARE_MAX/ 4.

no_of_bits_left_current_frame is the number of
bits left to send in the current frame at time CT. When
(latency < DTS_of_current_frame - CT), the numerator for
min_br represents the number of bits left to send in
the current frame at time CT+latency, and the
denominator represents the time period between
(CT+latency) and (DTS_of_current_frame). Min_br is
then the smallest bit rate that will enable the
remainder of the current frame to be transmitted when
15 implemented in the time period between (CT+latency) and
(DTS_of_current_frame). See FIG. 4(c).
20

```

Note that setting min\_br=0 essentially maintains the current transmission bit rate.

The code "min\_br = HARDWARE\_MAX/ 4" denotes a  
25 "panic mode" where there is not enough bandwidth to send the rest of the picture across the communication channel to the decoder before the DTS of the current picture. HARDWARE\_MAX is 28 Mbits/sec for SD and 39

Mbits/sec for HD. An alternative is to set `min_br = (no_of_bits_left_current_frame - (current_br - margin) * latency) / (one_QLBR_cycle)`. In this case, whatever is leftover in the current frame must be sent during the new QL/BR cycle.

Note that min and max bit rates are converted back into transport bit rates using a scale factor of 188/184 before they are sent back to the rate control processor, since the rate control processor deals with all transmission bit rates at the transport level.

1.3 A C-language syntax for a transmission bit rate process in accordance with the invention

(1) Initialize the bit rate parameters to nominal values.

```
15   for (i=0; i<num_mem; i++){
        min_br[i] = nom_br;
        max_br[i] = nom_br;
        target_br[i] = nom_br;
        bit_rate[i] = nom_br;
20      br_avail = br_avail - nom_br;
    }
```

(2) Calculate the total hard minimum transmission bit rate for the stat mux, then bound the minimum and maximum bit rates by the hard limits. Ignore the hard minimum if the total hard minimum exceeds the available bandwidth. Optionally, the user can set individual

hard limits for each channel.

```

tot_hmin = 0;

for (i=0; i<num_mem; i++)
    tot_hmin = tot_hmin + hminbr[i]; /* total hard minimum */

5      tot_hmin = 1.5 * tot_hmin;

for (i=0; i<num_mem; i++){
    if (min_br[i] < hminbr[i])
        if (tot_hmin < br_avail)
            min_br[i] = hminbr[i];

10     if (min_br[i] > hmaxbr[i])
            min_br[i] = hmaxbr[i];

        if (max_br[i] > hmaxbr[i])
            max_br[i] = hmaxbr[i];

        if (max_br[i] < hminbr[i])
15        if (tot_hmin < br_avail)
            max_br[i] = hminbr[i];
    }
}

```

(3) If the total minimum bit rate requested by the packet processor 250 (PP) of a channel encoder exceeds the available bandwidth, the rate control processor 125

handles the transmission bit rate allocation differently depending on whether or not the total hard minimum exceeds the available bandwidth. If it does, the bandwidth is allocated in proportion to the minimum bit rate requested by the packet processor. Otherwise, the rate control processor resorts to the hard minimums and allocates the remaining bandwidth in proportional to the difference between the requested minimums and the hard minimums.

```

10      total_min = tot_dmin = 0;

      for (i=0; i<num_mem; i++)
          tot_min = tot_min + min_br[i];

      if (tot_min > br_avail){ /* we don't have enough for all mins */
          for (i=0; i<num_mem; i++)
              tot_dmin = tot_dmin + (min_br[i] - hminbr[i]);
      }

      if (tot_hmin > br_avail){
          for (i=0; i<num_mem; i++) /* allocate proportional to minbr[i] */
              bit_rate[i] = br_avail * minbr[i] / tot_min;
      }
20      else {
          for (i = 0; i < num_mem; i++){
              bit_rate[i] = hminbr[i]; /* allocate hard minimums first */
              br_avail = br_avail - hminbr[i];
          }
      }
  
```

```

if (br_avail > 0){
    for (i = 0; i < num_mem; i++)
        bit_rate[i] = bit_rate[i] + br_avail * (min_br[i] - hminbr[i]) /
tot_dmin;                                /* use diff in requested and hard mins */

5          }
}
/* at this point no more br_avail */

br_avail = 0;
}

10      (4) If the total minimum bit rate requested by
the packet processor of a channel does not exceed the
available bandwidth, then allocate the minimum
requested bandwidth to each channel, and adjust the
target bit rate for each channel.

15      tot_target_br = 0;

for (i=0;i<num_mem;i++){
    br_avail=br_avail - min_br[i];
    bit_rate[i] = min_br[i];
    if (target_br[i] > min_br[i])
20        target_br[i] = target_br[i] - min_br[i];
    else
        target_br[i] = 0;
}

```

```

tot_target_br = tot_target_br + target_br[i];
}

```

(5) If there is still bandwidth available,  
distribute it in proportion to the channels' target  
5 bandwidths.

```

sav_bravail = br_avail;
if (br_avail > 0){
    for (i=0; i<num_mem; i++){
        if (tot_target_br != 0){
            bit_rate[i] = bit_rate[i] +
                (sav_bravail * target_br[i] / tot_target_br);
            br_avail = br_avail -
                (sav_bravail * target_br[i] / tot_target_br);
        }
        15    if (bit_rate[i] > max_br[i]){
            bit_rate[i] = max_br[i];
            br_avail = br_avail + (bit_rate[i] - max_br[i]);
        }
    }
    20}
}

```

(6) If there is still bandwidth available,  
distribute it in proportion to the difference between  
the max\_bitrate and the bit rate assigned, without  
exceeding the max\_bitrate.

```
if (br_avail > 0){  
    tot_excess_br = 0;  
    for (i=0; i<num_mem; i++){  
        if (bit_rate[i] < max_br[i])  
            excess_br[i] = max_br[i]-bit_rate[i];  
        else  
            excess_br[i] = 0;  
        tot_excess_br = tot_excess_br + excess_br[i];  
    }  
    if (tot_excess_br <= br_avail){  
        for(i=0; i<num_mem; i++)  
            bit_rate[i] = max_br[i];  
        br_avail = br_avail - tot_excess_br;  
    }  
    else{  
        for (i=0; i<num_mem; i++)  
            bit_rate[i] = bit_rate[i] + br_avail * excess_br[i]/tot_excess_br;  
    }  
}
```

FIG. 5 illustrates a method for setting minimum and maximum limits on a transmission bit rate in accordance with the present invention. A summary of the above discussions is provided.

An encoding bit rate is first allocated for a current picture (block 500), and a transmission bit rate is allocated in response thereto (block 510), and

following a system delay. If a new bit rate can be allocated before a decode time of the next picture, a max. bit rate limit is set to avoid a decoder buffer overflow at the decode time of the next picture (block 530). If a new bit rate can not be allocated before the decode time of the next picture, a max. bit rate limit is set to avoid a decoder buffer overflow at the decode time of the next, next picture (block 540).

For setting a minimum limit on the transmission bit rate (550), e.g., to avoiding a possible decoder buffer underflow, if a new bit rate can be allocated before the decode time of the current picture, the min. bit rate is set so that a remainder of the current picture is transmitted before the decode time of the current picture (block 560). If a new bit rate can not be allocated before the decode time of the current picture, a determination is made as to whether the current transmission bit rate is sufficient to transmit the remainder of the current picture before the decode time of the current picture (block 570). If the current bit rate is sufficient, the min. bit rate is set to zero (block 580) since there is no need to force the bit rate higher. If the current bit rate is not sufficient, the min. bit rate is set to a maximum value, such as a hardware maximum (block 590), to force the bit rate to the highest possible level to mitigate any possible buffer underflow at the decode time of the current picture.

The transmission bit rate for the current picture is adjusted as necessary in view of the min. and max. limits (block 595). The processing then repeats for the next transmission bit rate allocation cycle. The min. and max. limits take into account a delay in implementing a new bit rate, and also account for whether the new bit rate can be implemented before the decoding time of the current picture, the next picture, and/or the next, next picture.

Accordingly, it can be seen that the present invention provides a statistical multiplexer for coding and multiplexing multiple channels of digital television data. A bit rate need parameter is determined for each encoder in a stat mux group, and an encoding bit rate is allocated to each channel based on its need parameter.

Additionally, a transmission bit rate is allocated to each channel as a time-lagged version of its need parameter to minimize a rate mismatch between the output and the input of a decoder buffer. A packet processor checks for impending decoder buffer overflow or underflow events to set minimum and maximum limits on the transmission bit rate.

Although the invention has been described in connection with various preferred embodiments, it should be appreciated that various modifications and adaptations may be made thereto without departing from the scope of the invention as set forth in the claims.

What is claimed is:

1. A method for processing a plurality of channels in a statistical multiplexer, comprising the steps of:

allocating an encoding bit rate for coding a current picture of each channel according to a bit rate need parameter thereof; and

for each channel, allocating a transmission bit rate for transmitting the current picture after encoding thereof, and providing a modeled decoder buffer that receives transmitted pictures therefrom;

wherein, for each channel, the transmission bit rate is based on the channel's encoding bit rate, and is allocated following a system delay that follows the allocated encoding bit rate, to minimize a rate mismatch between an input and an output of the modeled decoder buffer.

2. The method of claim 1, wherein:

for at least one of the channels, the transmission bit rate is allowed to deviate from the channel's encoding bit rate to avoid an impending overflow or underflow event of the modeled decoder buffer.

3. The method of claim 1, comprising the further step of:

for at least one of the channels, checking for

impending overflow or underflow events of the modeled decoder buffer to set at least one of minimum and maximum limits on the transmission bit rates for the channel.

4. The method of claim 1, wherein  
for at least one of the channels, when a next update of the allocated transmission bit rate can be implemented, following an implementation delay, before a decode time of a next picture, a maximum limit is set on the allocated transmission bit rate at a current time (CT) in proportion to a fullness of the modeled decoder buffer at a time (CT+delay), and in inverse proportion to a time period between (CT+delay) and the decode time.

5. The method of claim 1, wherein  
for at least one of the channels, when a next update of the allocated transmission bit rate can not be implemented, following an implementation delay, before a decode time of a next picture, a maximum limit is set on the allocated transmission bit rate at a current time (CT) in proportion to a fullness of the modeled decoder buffer at a time (CT+delay), and in inverse proportion to a time period between (CT+delay) and a decode time of a picture that follows said next picture.

6. The method of claim 1, wherein  
for at least one of the channels, when a next  
update of the allocated transmission bit rate can be  
implemented, following an implementation delay, before  
a decode time of the current picture, a minimum limit  
is set on the allocated transmission bit rate at a  
current time (CT) in proportion to a number of  
remaining bits of the current picture to transmit at a  
time (CT+delay), and in inverse proportion to a time  
period between (CT+delay) and the decode time.

7. The method of claim 1, comprising the further  
step of:

for at least one of the channels, determining  
whether a current allocated transmission bit rate is  
sufficient to transmit a number of remaining bits of  
the current picture in a time period between a current  
time and a decode time of the current picture, and, if  
so, maintaining the current allocated transmission bit  
rate in a next update cycle thereof.

8. The method of claim 1, comprising the further  
step of:

for at least one of the channels, forcing the  
allocated transmission bit rate to a maximum value in a  
next update cycle thereof when a current allocated  
transmission bit rate is not sufficient to transmit a  
number of remaining bits of the current picture in a

time period between a current time and a decode time of the current picture.

9. The method of claim 1, wherein  
for at least one of the channels, when a next update of the allocated transmission bit rate can be implemented, following an implementation delay, before a decode time of a next picture, a maximum limit is set on the allocated transmission bit rate at a current time to avoid an overflow of the modeled decoder buffer at the decode time.

10. The method of claim 1, wherein  
for at least one of the channels, when a next update of the allocated transmission bit rate can not be implemented, following an implementation delay, before a decode time of a next picture, a maximum limit is set on the allocated transmission bit rate at a current time to avoid an overflow of the modeled decoder buffer at a decode time of a picture that follows said next picture.

11. The method of claim 1, wherein:  
for at least one of the channels, when a next update of the allocated transmission bit rate can be implemented, following an implementation delay, before a decode time of the current picture, a minimum limit is set on the allocated transmission bit rate at a

current time such that the current picture is completely transmitted before the decode time.

12. The method of claim 1, wherein:

for at least one of the channels, when a next update of the allocated transmission bit rate can not be implemented, following an implementation delay, before a decode time of the current picture, a minimum limit on the allocated transmission bit rate at a current time is set to a maximum value to mitigate a potential underflow of the modeled decoder buffer.

~~13.~~ An apparatus for processing a plurality of channels in a statistical multiplexer, comprising:

means for allocating an encoding bit rate for coding a current picture of each channel according to a bit rate need parameter thereof; and

means for allocating a transmission bit rate for transmitting the current picture after encoding thereof, and providing a modeled decoder buffer that receives transmitted pictures therefrom, for each channel;

wherein, for each channel, the transmission bit rate is based on the channel's encoding bit rate, and is allocated following a system delay that follows the allocated encoding bit rate, to minimize a rate mismatch between an input and an output of the modeled decoder buffer.

**ABSTRACT OF THE DISCLOSURE**

A statistical multiplexer for coding and multiplexing multiple channels of digital television data, or multiple panels of HDTV digital television data. A bit rate need parameter is determined for each encoder in a stat mux group, and an encoding bit rate is allocated to each channel based on its need parameter. A transmission bit rate is allocated to each channel as a time-lagged version of its need parameter to minimize a rate mismatch between the output and the input of a decoder buffer. A packet processor checks for impending decoder buffer overflow or underflow events to set minimum and maximum limits on the transmission bit rate. Moreover, these limits are set based on whether a new transmission bit rate can be implemented before the decoding time stamp (DTS) of the current or next frame.

P A T E N T

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Application of:

)

Liu et al.

)

Application No.:

)

Filed: Herewith

)

)

For: **METHOD AND APPARATUS FOR DETERMINING A TRANSMISSION BIT RATE  
IN A STATISTICAL MULTIPLEXER**

**DRAWING REVIEW BRANCH**

Commissioner for Patents  
Washington, D.C. 20231

CERTIFICATE OF MAILING

I hereby certify that this correspondence is being deposited with the United States Postal Service with sufficient postage as Express Mail (No EL 632260048 US) addressed to BOX PATENT APPLICATION, Commissioner for Patents, Washington, D C 20231 on September 20, 2000.

By Michele Holis  
Michele Holis

TRANSMITTAL OF FORMAL DRAWINGS

Dear Sir:

Enclosed are seven (7) sheets of formal drawings for filing in the above-referenced patent application.

Please advise the undersigned attorney if correction is necessary.

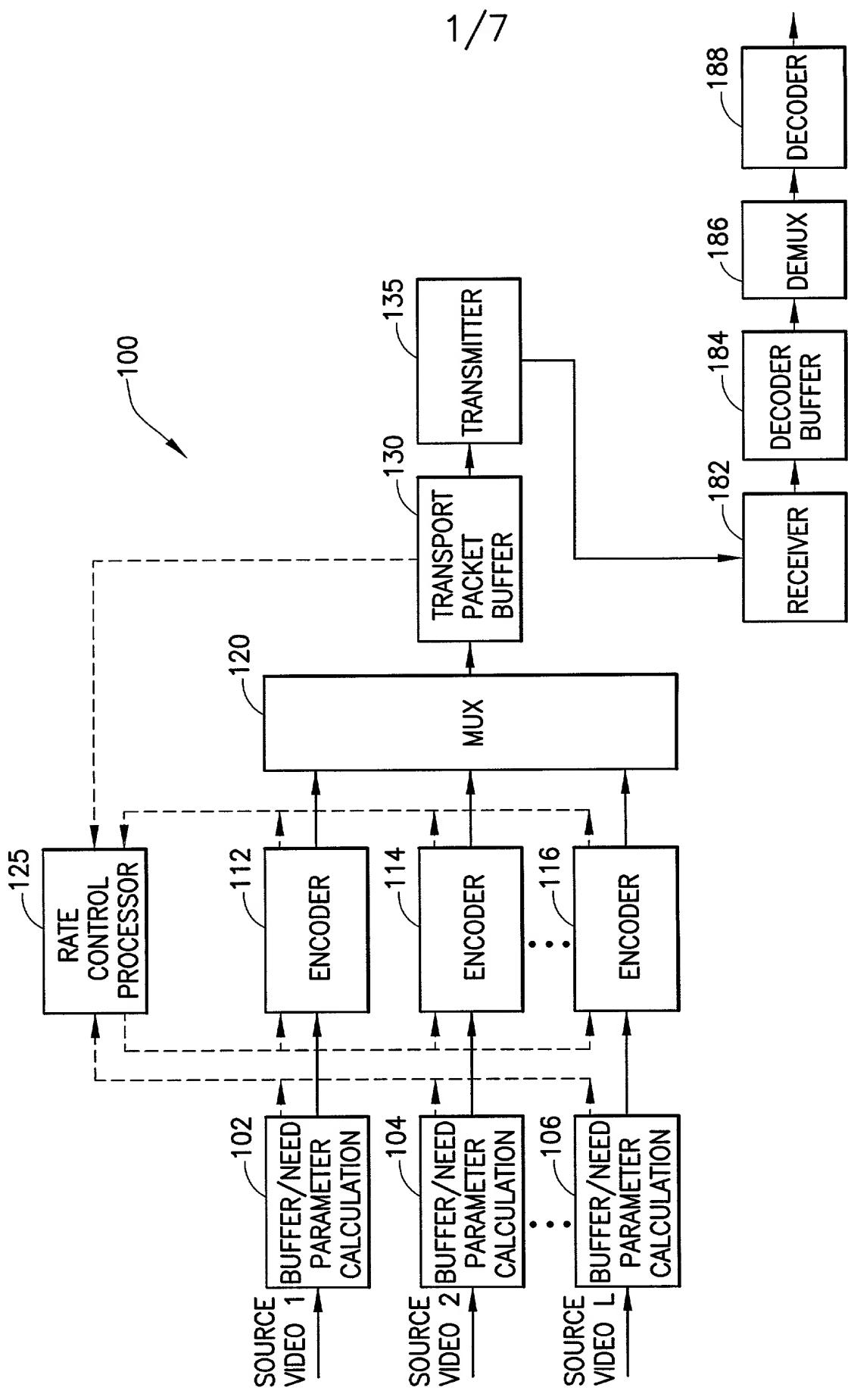
Respectfully submitted,



Barry R. Lipsitz  
Attorney for Applicant(s)  
Registration No. 28,637  
755 Main Street  
Monroe, CT 06468  
(203) 459-0200

Date: September 20, 2000  
**ATTORNEY DOCKET NO.: GIC-620**

FIG. 1



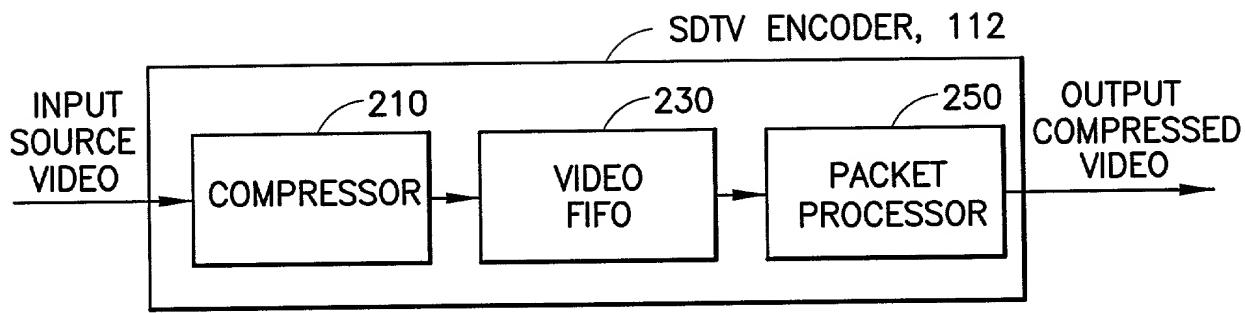
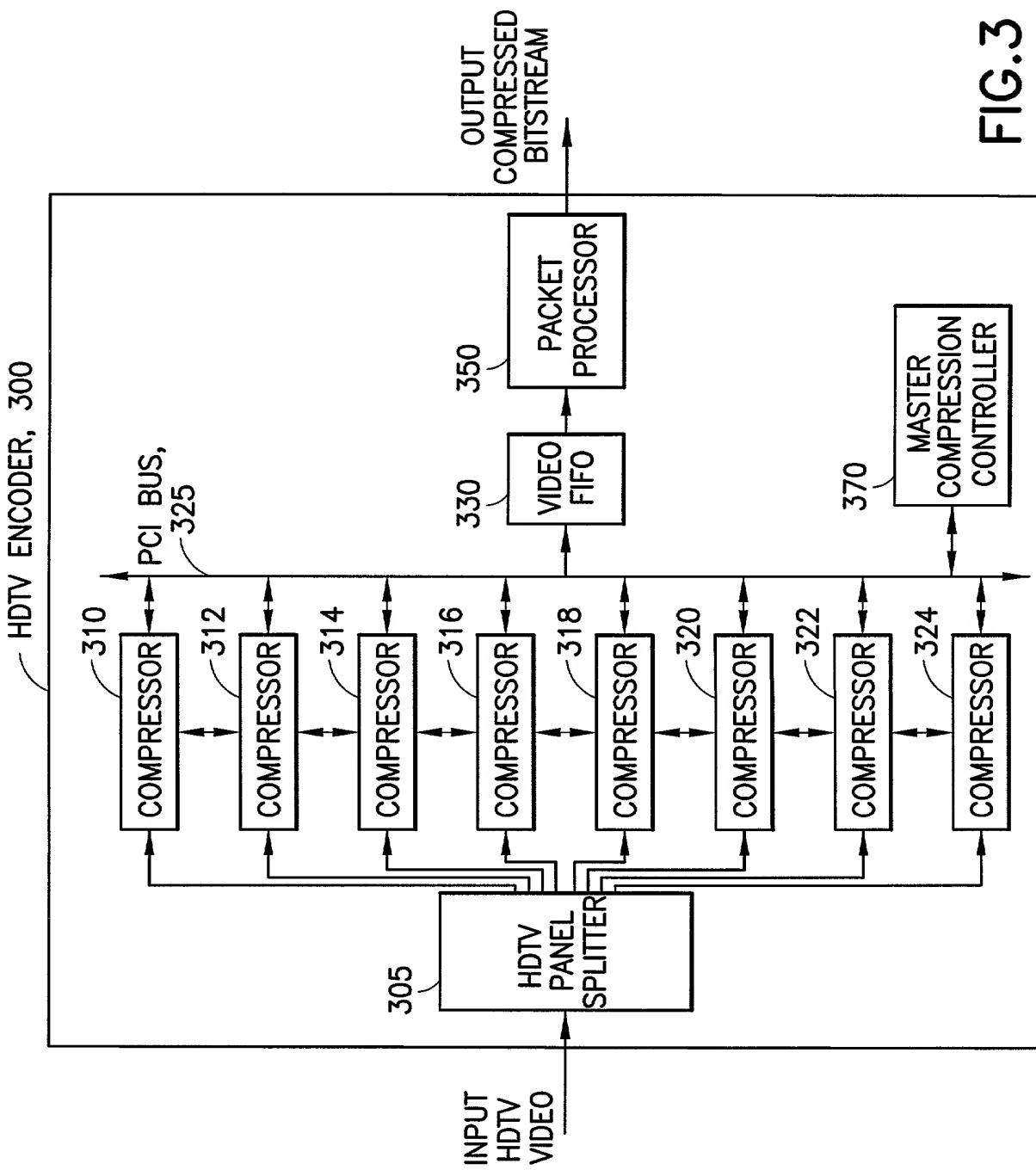


FIG.2



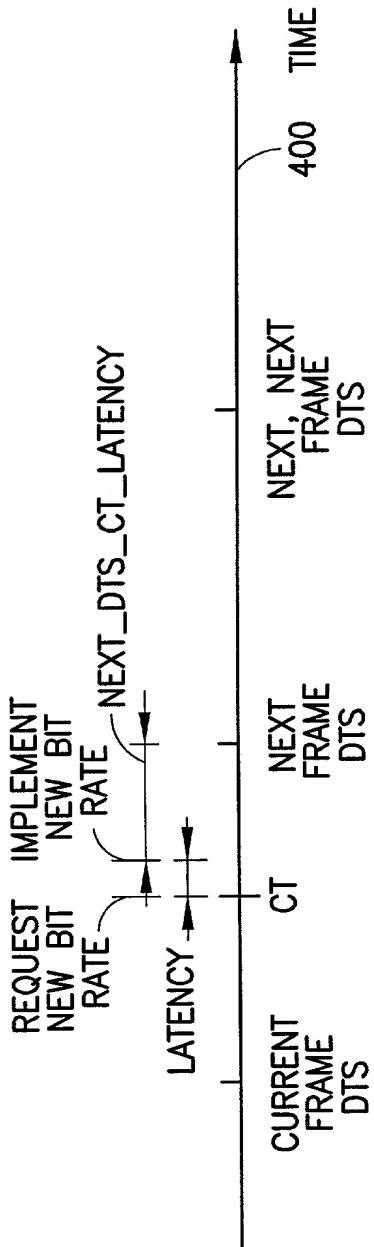


FIG.4(a)

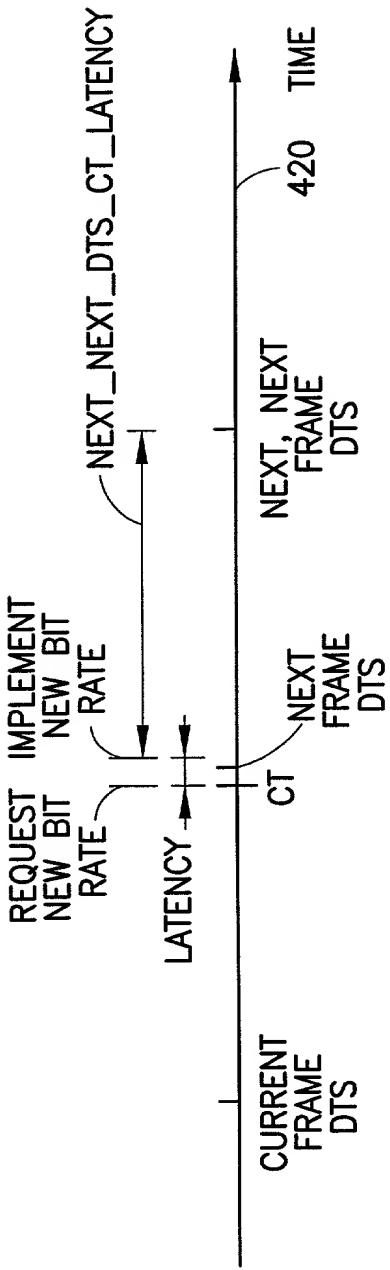


FIG.4(b)

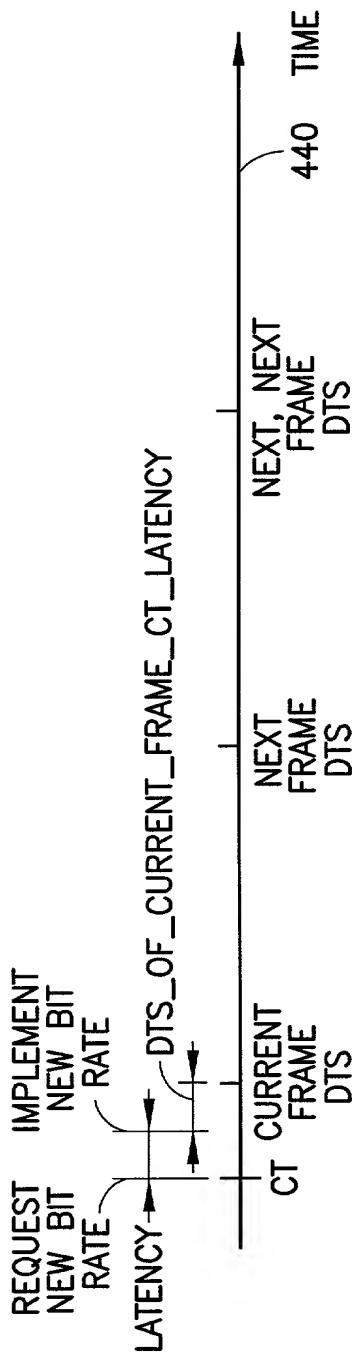


FIG.4(c)

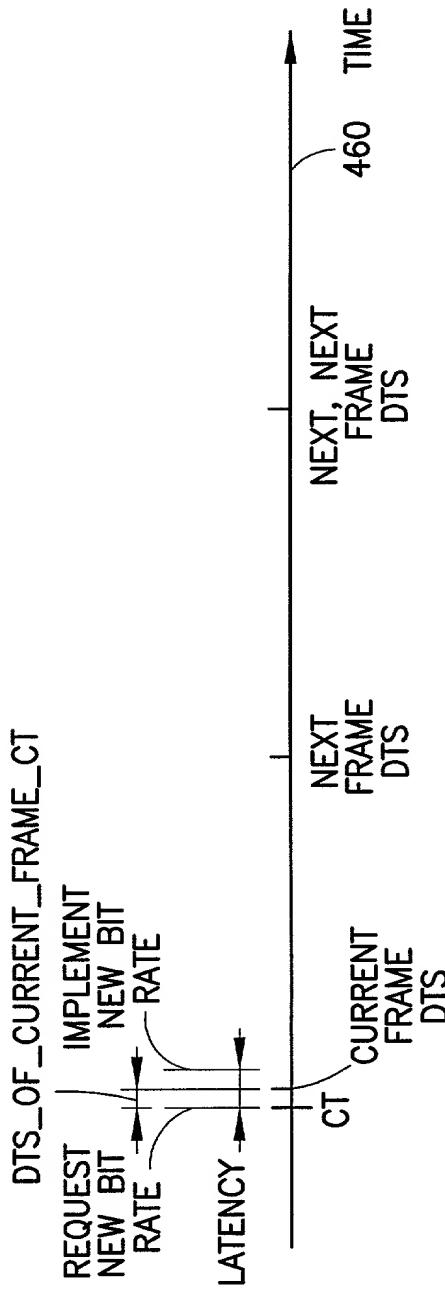


FIG.4(d)

500

ALLOCATE AN ENCODING BIT RATE  
FOR A CURRENT PICTURE BASED  
ON ITS NEED PARAMETER

FIG. 5A

FIG. 5B

FIG. 5

510

ALLOCATE A TRANS. BIT RATE  
BASED ON THE ENCODING BIT RATE,  
AND FOLLOWING A SYSTEM DELAY

6 / 7

FIG.5A

FIG.5B

520

CAN A NEW TRANS. BIT RATE BE  
IMPLEMENTED BEFORE A DECODE  
TIME OF A NEXT PICTURE?

YES

530

SET A MAX TRANS. BIT RATE  
TO AVOID A MODELED  
DECODER BUFFER OVERFLOW  
AT THE DECODE TIME OF  
THE NEXT PICTURE

540

SET A MAX TRANS. BIT RATE  
TO AVOID A MODELED  
DECODER BUFFER OVERFLOW  
AT A DECODE TIME OF THE  
NEXT, NEXT PICTURE

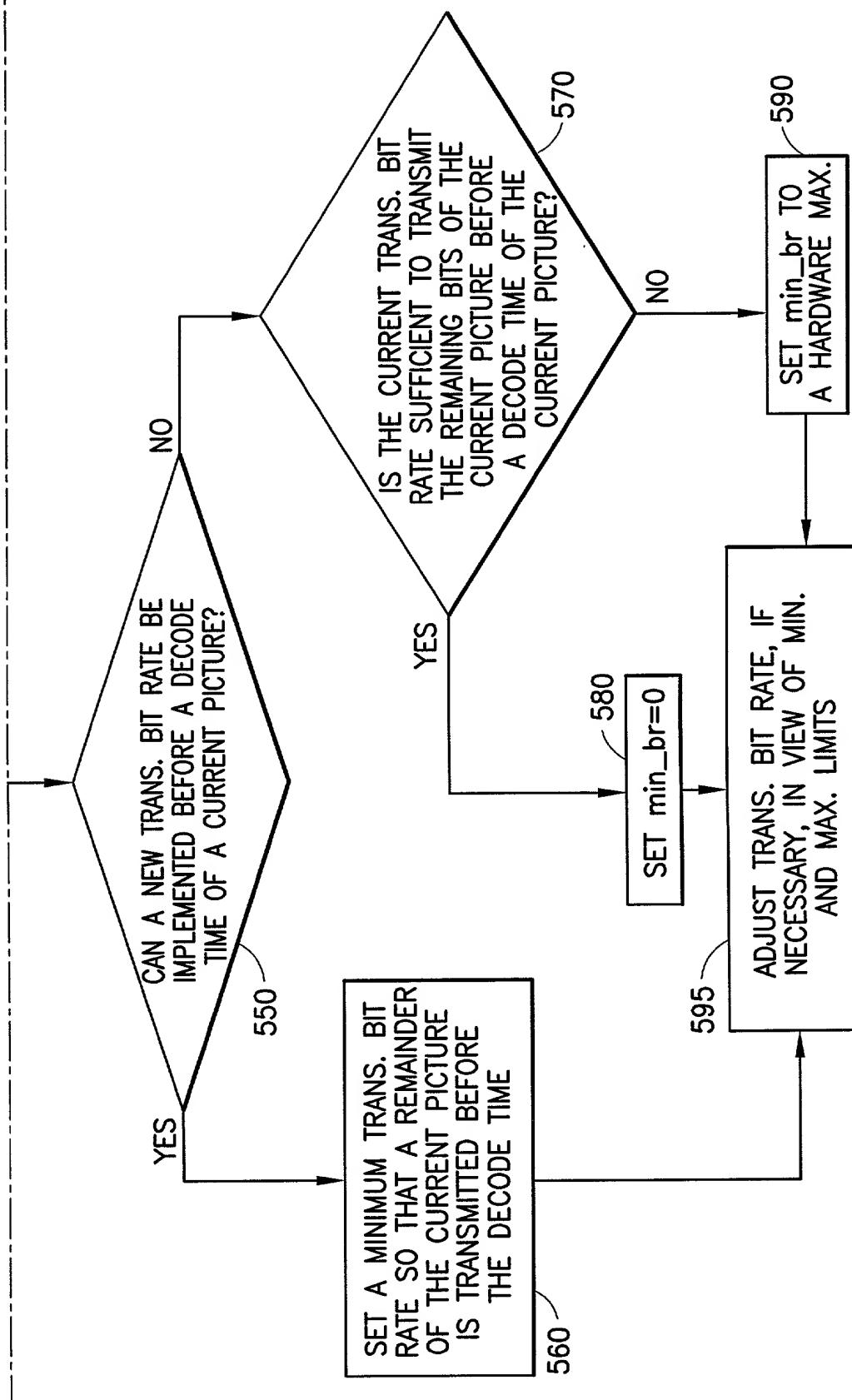


FIG.5B

## **DECLARATION, POWER OF ATTORNEY, AND PETITION**

Attorney Docket No.: GIC-620  
Page 1 of 3

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled

### **METHOD AND APPARATUS FOR DETERMINING A TRANSMISSION BIT RATE IN A STATISTICAL MULTIPLEXER**

the specification of which is attached hereto unless the following box is checked:

[ ] was filed on \_\_\_\_\_ as United States Application Number \_\_\_\_\_ or PCT International Application Number \_\_\_\_\_ and was amended on \_\_\_\_\_ (if applicable).

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose to the U.S. Patent and Trademark Office all information known to be material to the patentability of this application in accordance with Title 37, Code of Federal Regulations, §1.56.

I hereby claim foreign priority benefits under Title 35, United States Code, §119(a)-(d) or 365(b) of any foreign application(s) for patent or inventor's certificate or 365(a) of any PCT international application which designated at least one country other than the United States of America, listed below and have also identified below any foreign application for patent or inventor's certificate or of any PCT international application having a filing date before that of the application on which priority is claimed:

			Priority Claimed
(Number)	(Country)	Month/Day/Year Filed	[ ] [ ] Yes No
(Number)	(Country)	Month/Day/Year Filed	[ ] [ ] Yes No

I hereby claim the benefit under Title 35, United States Code, §119(e) of any United States provisional application(s) listed below.

(Application Number) \_\_\_\_\_ (Filing Date) - Month/Day/Year \_\_\_\_\_

(Application Number) \_\_\_\_\_ (Filing Date) - Month/Day/Year \_\_\_\_\_

I hereby claim the benefit under 35 U.S.C. 120 of any United States application(s), or 365(c) of any PCT international application designating the United States of America, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT international application in the manner provided by the first paragraph of 35 U.S.C. 112, I acknowledge the duty to disclose information which is material to patentability as defined in 37 C.F.R. 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application.

**U.S. Parent Application  
or PCT Parent Number**

**Parent Filing Date  
(MM/DD/YYYY)**

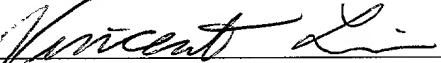
**Parent Patent Number  
(if applicable)**

And I hereby appoint: Barry R. Lipsitz, Registration No. 28,637, Ralph F. Hoppin, Registration No. 38,494 and Douglas M. McAllister, Registration No. 37,886, all of the firm of Barry R. Lipsitz, Attorney at Law, 755 Main Street, Bldg. 8, Monroe, Connecticut 06468, Telephone (203) 459-0200, my attorneys with full power of substitution and revocation, to prosecute this application and to transact all business in the Patent and Trademark Office connected therewith.

Wherefore I pray that Letters Patent be granted to me for the invention or discovery described and claimed in the foregoing specification and claims, and I hereby subscribe my name to the foregoing specification and claims, declaration, power of attorney, and this petition.

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

**Full name of sole or first inventor:** Vincent LIU

Inventor's Signature  Date: 8-24-2000

Residence San Diego California TAIWAN  
(City) (State or Foreign Country) Citizenship:  
  
Post Office Address 7714 Roan Road San Diego, California 92129, U.S.A.  
(Post Office Address) (City) (State & Zip Code/Country)

**Full name of second joint inventor:** Siu-Wai WU

Inventor's Signature  Date: 8-24-2000

Residence San Diego California HONG KONG  
(City) (State or Foreign Country) Citizenship:  
  
Post Office Address 12655 Fairbrook Road San Diego, California 92131, U.S.A.  
(Post Office Address) (City) (State & Zip Code/Country)

**Full name of third joint inventor:** Michael CASTELOES

Inventor's Signature  Date: 8-31-00

Residence Poway California U.S.A.  
(City) (State or Foreign Country) Citizenship:  
  
Post Office Address 14227 Woodcreek Road Poway, California 92064, U.S.A.  
(Post Office Address) (City) (State & Zip Code/Country)

Full name of fourth joint inventor: Robert J. STONE

Inventor's Signature

Date: 8/31/00

Residence Poway California United Kingdom  
(City) (State or Foreign Country) Citizenship:

Post Office Address 13558 Quiet Hills Drive Poway, California 92064, U.S.A.  
(Post Office Address) (City) (State & Zip Code/Country)

Full name of fifth joint inventor: Rebecca LAM

Inventor's Signature

Date: 8/31/2000

Residence San Diego California U.S.A.  
(City) (State or Foreign Country) Citizenship:

Post Office Address 11933 Acacia Glen Ct. San Diego, California 92128, U.S.A.  
(Post Office Address) (City) (State & Zip Code/Country)